The Origins of the Theory of Random Graphs

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1. Introduction

The origins of the theory of random graphs are easy to pin down. Undoubtedly 15 one should look at a sequence of eight papers co-authored by two great 16 mathematicians: Paul Erdős and Alfred Rényi, published between 1959 and 17 1968: 18

[ER59] On random graphs I, Publ. Math. Debrecen 6 (1959), 290–297.

- [ER60] On the evolution of random graphs, Publ. Math. Inst. Hung. Acad. 20
 Sci. 5 (1960), 17–61.
- [ER61a] On the evolution of random graphs, Bull. Inst. Internat. Statist. 38, 22 343–347.
- [ER61b] On the strength of connectedness of a random graph, Acta Math. 24
 Acad. Sci. Hungar. 12 (1961), 261–267.

[ER63] Asymmetric graphs, Acta Math. Acad. Sci. Hung. 14, 295–315.

[ER64] On random matrices, Publ. Math. Inst. Hung. Acad. Sci. 8 (1964), 27 455–461.

[ER66] On the existence of a factor of degree one of a connected random 29 graph, Acta Math. Acad. Sci. Hung. **17** (1986), 359–368.

[ER68] On random matrices II, Studia Sci. Math. Hung. 3 (1968), 459–464. 31

Our main goal is to summarize the results, ideas and open problems 32 contained in those contributions and to show how they influenced future 33 research in random graphs. 34

For us it was a great adventure to return to the roots of the theory of 35 random graphs, and to find out again and again, how far-reaching the impact 36 of Erdős and Rényi's work on the field is. The reader will find in our paper 37 many quotations from their original papers (*always in italics*). We use this 38 convention to let them speak directly and to preserve their special insightful 39

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style and way of thinking and stating the problems. Starting from there we 40 lead the reader through the literature, including the most current one, trying 41 to show how the ideas of Erdős and Rényi developed, how much time, skills 42 and effort to solve some of their most challenging open problems was needed. 43 Finally, to add some "salt and pepper" to our presentation, full of admiration 44 and respect, we point out to a few false statements and oversimplifications 45 of proofs, which have been found in their monumental legacy by the next 46 generations of random graph theorists.

2. The First Question: Connectivity

Although the notion of a random graph appeared in connection to the 49 probabilistic method already in the Erdős paper [25] (see J. Spencer's article 50 in this volume), it was forgotten for a decade until Paul Erdős and Alfred 51 Rényi published a series of papers entirely devoted to properties of random 52 graphs. The model of a random graph they exclusively investigated was the 53 uniform one. Here is how they defined it: "Let $E_{n,N}$ denote the set of all 54 graphs having n given labeled vertices and N edges. A random graph $\Gamma_{n,N}$ 55 can be defined as an element of $E_{n,N}$ chosen at random, so that each of the 56 elements of $E_{n,N}$ have the same probability to be chosen, namely $1/{\binom{n}{2}}{N}$." 57 (In this paper we adopt the original notation $\Gamma_{n,N}$.)

They were aware of existing results about other models of random graphs. ⁵⁹ In particular, they acknowledge in a footnote to [ER61a] that E. N. Gilbert ⁶⁰ [36] studied the connectedness of what we call today the binomial model, ⁶¹ where "We may decide with respect to each of the $\binom{n}{2}$ edges, whether they ⁶² should form part of the random graph considered or not, the probability of ⁶³ including a given edge being $p = N/\binom{n}{2}$ for each edge and the decisions ⁶⁴ concerning different edges being independent." (In this paper we shall denote ⁶⁵ this model by $\Gamma_{n,p}$.) In [ER61a] they mention that the investigations of ⁶⁶ the binomial model can be reduced, due to a conditional argument they ⁶⁷ attribute to Hajek, to that of $\Gamma_{n,N}$. However, they did not formulate ⁶⁸ any equivalence theorem (these appeared much later in [14] and [59]) and ⁶⁹ occasionally stated the binomial counterparts of their theorems without ⁷⁰ proofs or repeated their proofs step by step. ⁷¹

Apparently they were not aware of the result of Gilbert and of the binomial model at all when they wrote their first paper on random graphs, "On 73 random graphs I". The question addressed there was that of connectedness of 74 a random graph. In fact, according to a remark in [ER59], this problem was 75 tried and partially solved already in 1939, when P. Erdős and H. Whitney, 76 in an unpublished work: "proved that if $N > (\frac{1}{2} + \varepsilon) n \log n$ where $\varepsilon > 0$ 77 then the probability of $\Gamma_{n,N}$ being connected tends to 1 if $n \to \infty$, but if 78 $N < (\frac{1}{2} - \varepsilon) n \log n$ with $\varepsilon > 0$ then the probability of $\Gamma_{n,N}$ being connected, 79 tends to 0 if $n \to \infty$."

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In the first "official" paper on random graphs, Erdős and Rényi refined 81 the above result as their (partial) answer to questions 1–3 from the following 82 list of problems they posed. 83

- 1. What is the probability of $\Gamma_{n,N}$ being completely connected?
- 2. What is the probability that the greatest connected component (subgraph) 85 of $\Gamma_{n,N}$ should have effectively n-k points? $(k=0,1,\ldots)$ 86
- 3. What is the probability that $\Gamma_{n,N}$ should consist of exactly k+1 connected so components? (k = 0, 1, ...)
- 4. If the edges of a graph with n vertices are chosen successively so that after s9 each step every edge which has not yet been chosen has the same probability 90 to be chosen as the next, and if we continue this process until the graph 91 becomes completely connected, what is the probability that the number of 92 necessary steps ν will be equal to a given number l?

Note that in problem 4 Erdős and Rényi describe a genuine random graph 94 process, whose advanced analysis could be carried over only two decades later. 95

Before turning to the proofs, they recall a recursive formula and a 96 generating function for the number C(n, N) of connected graphs on n labeled 97 vertices and with N edges, due to Riddell and Uhlenbeck, and also Gilbert. 98 But immediately they comment that neither of them "...helps much to 99 deduce the asymptotic properties of C(n, N). In the present paper we follow 100 a more direct approach." 101

We now present the first result on random graphs and its proof in a 102 slightly modified form. The idea of the proof, however, remains unchanged. 103 In the 1959 paper only the middle part of the theorem below was stated 104 explicitly. The other two follow by letting $c = c_n$ tend to $+\infty$ or $-\infty$, 105 respectively. 106

Theorem 2.1 ([26])

$$P(\Gamma_{n,N} \text{ is connected }) \to \begin{cases} 0 & \text{if } \frac{N}{n} - \frac{1}{2}\log n \to -\infty \\ e^{-e^{-2c}} & \text{if } \frac{N}{n} - \frac{1}{2}\log n \to c \\ 1 & \text{if } \frac{N}{n} - \frac{1}{2}\log n \to \infty. \end{cases}$$

Proof. For convenience we switch to the binomial model, shortening the 109 original argument a lot, and, at the same time, avoiding a harmless error 110 in the proof of "the rather surprising Lemma" of [ER59], pointed out by 111 Godehardt and Steinbach [37].

To make this argument formal, assume that $2np - \log n - \log \log n \to \infty$ 113 but $np = O(\log n)$. Thus, almost surely (i.e., with probability tending to 1 114 as $n \to \infty$), there are no isolated edges in $\Gamma_{n,p}$. What remains to be shown 115 is that there are no components of size $3 \le k \le \frac{n}{2}$ either. To this end 116 consider the random variable X counting such components. Then, bounding

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the probability that a given set of k vertices spans a connected subgraph by 117 $k^{k-2}p^{k-1},$ and using the inequality $np>\frac{1}{2}\log n,$ we obtain 118

$$Exp(X) \le \sum_{k=3}^{n/2} \binom{n}{k} k^{k-2} p^{k-1} (1-p)^{k(n-k)} < \sum_{k} \left(\frac{en}{k}\right)^{k} k^{k-2} p^{k-1} e^{-(n-k)pk}$$
$$\le \frac{1}{p} \sum_{k=3}^{\sqrt{n}} \frac{1}{k^{2}} \left(\frac{enp}{e^{(n-\sqrt{n})p}}\right)^{k} + \frac{1}{p} \sum_{k\ge\sqrt{n}}^{n} \frac{1}{n} \left(\frac{enp}{e^{np/2}}\right)^{k}$$
$$= O\left(\frac{n}{\log n} \frac{\log^{3} n}{n^{3/2}}\right) + \frac{1}{\log n} \left(\frac{e\log n}{2n^{1/4}}\right)^{\sqrt{n}} = o(1).$$

Hence, almost surely there are no components outside the largest 119 one other than isolated vertices (Erdős and Rényi say that such a graph 120 is of type A) and the threshold for connectedness coincides with that for 121 disappearance of isolated vertices, i.e., for $2np - \log n - \log \log n \rightarrow \infty$ 122

$$P(\Gamma_{n,p} \text{ is connected }) = P(\delta(\Gamma_{n,p}) > 0) + o(1).$$
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Erdős and Rényi found the limiting value of $P(\delta(\Gamma_{n,p}) > 0)$ by inclusionexclusion. Nowadays a standard approach is by the method of moments which serves to show that the number of isolates is asymptotically Poisson. They used that method in the 1960 paper in a more general setting where components isomorphic to a given graph G were considered. We shall return to this later.

Answering question 4, they gave a somewhat oversimplified proof of the 130 fact that 131

$$\lim_{n \to \infty} P\left(\frac{\nu - \frac{1}{2}n\log n}{n} < x\right) = e^{-e^{-2x}} . \qquad \Box \quad 132$$

Erdős and Rényi conclude the 1959 paper as follows. "The following more 134 general question can be asked: Consider the random graph $\Gamma_{n,N(n)}$ with n 135 possible vertices and N(n) edges. What is the distribution of the number of 136 vertices of the greatest connected component of $\Gamma_{n,N(n)}$ and the distribution of 137 the number of its components? What is the typical structure of $\Gamma_{n,N(n)}$ (in the 138 sense in which, according to our Lemma, the typical structure of $\Gamma_{n,N(n)}$ is 139 that it belongs to type A)? We have solved these problems in the present paper 140 only in the case $N(n) = \frac{1}{2}n\log n + cn$. We shall return to the general case in 141 another paper [8]." ([8] = [ER60] on our reference list.)

As far as connectedness is concerned, in the 1961 paper Erdős and 143 Rényi go on and find the threshold for r-connectivity of $\Gamma_{n,p}$ for every 144 natural r. "If G is an arbitrary non-complete graph, let $c_p(G)$ denote the 145 least number k such that by deleting k appropriately chosen vertices from 146 G (...) the resulting graph is not connected. (...) Let $c_e(G)$ denote the 147

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least number l such that by deleting l appropriately chosen edges from G 148 the resulting graph is not connected." A graph is r-connected if no removal 149 of r or less vertices can disconnect it. When the random graph becomes 150 almost surely r-connected? Theorem 2.1 revealed an interesting feature of 151 random graphs. Namely, quite often trivial necessary conditions become 152 asymptotically sufficient in the sense that for a typical, large graph their 153 fulfillment guaranties that the property in question holds. Due to Theorem 2.1 154 this is the case of connectedness versus the nonexistence of isolated vertices. 155 For r-connectedness such natural necessary condition is that the minimum 156 degree (denoted in [ER61b] by c(G)) must be at least r. Otherwise removing 157 the vertices adjacent to a vertex of minimum degree would disconnect the 158 graph. Erdős and Rényi showed in 1961 that in the range $\frac{1}{2}n\log n \leq N \leq$ 159 $n \log n$ this is the only way one can disconnect the random graph $\Gamma_{n,N}$ by 160 removing the smallest possible number of vertices. A minimal cutset is a set 161 of vertices whose removal makes the graph disconnected but no proper subset 162 of that set has this property. For $2 \le k \le \frac{n-1}{2}$ let \mathcal{A}_k be the event that there 163 is in $\Gamma_{n,N}$ a minimal cutset of size $s, 1 \leq s \leq r-1$, which leaves the second 164 largest component of size k. Arguing similarly as in the proof of Theorem 2.1, 165 they proved that $P(\bigcup_{k\geq 2} \mathcal{A}_k) = o(1)$, meaning that, almost surely, if $\Gamma_{n,N}$ 166 is not r-connected then the only reason for that is the presence of vertices of 167degree less than r. The method of moments (again, in the inclusion-exclusion 168 cover-up) gives that, for $N(n) = \frac{1}{2}n\log n + \frac{r}{2}n\log\log n + an + o(n)$, their 169 number is asymptotically Poisson. We thus arrived at the main result of the 170 1961 paper. 171

Theorem 2.2 ([28]). If we have

$$N(n) = \frac{1}{2}n\log n + \frac{r}{2}n\log\log n + an + o(n)$$
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where a is a real constant and r a non-negative integer, then 174

$$\lim_{n \to \infty} P(c_p(\Gamma_{n,N(n)}) = r) = 1 - \exp\left(-\frac{e^{-2a}}{r!}\right),\tag{3}$$

$$\lim_{n \to \infty} P(c_e(\Gamma_{n,N(n)}) = r) = 1 - \exp\left(-\frac{e^{-2a}}{r!}\right) \tag{4}$$

and

$$\lim_{n \to \infty} P(c(\Gamma_{n,N(n)}) = r) = 1 - \exp\left(-\frac{e^{-2a}}{r!}\right).$$
(5)

In a proceeding remark they promise: "The statement (5) of Theorem 2.2 177 gives information about the minimal valency of points of $\Gamma_{n,N}$. In a forth-178

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coming note we shall deal with the same question for larger ranges of N ¹⁷⁹ (when $c(\Gamma_{n,N})$ tends to infinity with n), further with the related question about ¹⁸⁰ maximal valency of points of $\Gamma_{n,N}$." This promise was never fulfilled. The only ¹⁸¹ trace of their interest in the vertex degrees of a random graph can be found in ¹⁸² the description of the last phase of the evolution of $\Gamma_{n,N}$ in [ER61a]: "Phase 5. ¹⁸³ consists of the range $N(n) \sim (n \log n)w(n)$ where $w(n) \to \infty$. In this range ¹⁸⁴ the whole graph is not only almost surely connected, but the orders of points ¹⁸⁵ are almost surely asymptotically equal. Thus the graph becomes in this phase ¹⁸⁶ 'asymptotically regular'." The proof of that statement can be found in the ¹⁸⁷ last section of [ER60]. A very careful analysis of vertex degrees in a random ¹⁸⁸ graph is due to Bollobás [10, 11] and can be found also in his book [14].

3. Subgraphs: The Beginning of a Theory

After having written their paper on connectivity of a random graph Erdős and 191 Rényi decide to write a long paper addressing several properties of random 192 graphs. That seminal paper was preceded by an extended abstract [ER61a], 193 where they outlined the main goals of the theory to be born. "Our main goal 194 is to show (...) that the evolution of a random graph shows very clear-cut 195 features. The theorems we have proved belong to two classes. The theorems of 196 the first class deal with the appearance of certain subgraphs (e.g., tress, cycles 197 of a given order etc.) or components, or other local structural properties, and 198 show that for many types of local structural properties A a definite 'threshold' 199 A(n) can be given, so that if $\frac{N(n)}{A(n)} \to 0$ for $n \to \infty$ then the probability 200 that the random graph $\Gamma_{n,N(n)}$ has the structural property A tends to 0 for 201 $n \to \infty$, while for $\frac{N(n)}{A(n)} \to \infty$ for $n \to \infty$ the probability that the random 202 graph $\Gamma_{n,N(n)}$ has the structural property A tends to 1 for $n \to \infty$. (...) 203 The theorems of the second class are of similar type, only the properties A 204 considered are not of a local character, but global properties of the graph 205 $\Gamma_{n,N(n)}$ (e.g., connectivity, total number of components, etc.)." The existence 206 of a threshold in all cases they considered was a rather surprising fact for 207 Erdős and Rényi. Only three decades later it was proved by Bollobás and 208 Thomason [19] that, as a consequence of the Kruskal-Katona inequality, every 209 monotone property (family) of random subsets of a set has a threshold in the 210 above sense. 211

In the same abstract they comment that their proofs are "... completely 212 elementary, and are based on the asymptotic evaluation of combinatorial 213 formulae and on some well-known general methods of probability theory" 214

The first theorem of the major paper [ER60] established the threshold for 215 the existence of a subgraph of a given type for a broad class of subgraphs. 216 "If a graph has n vertices and N edges, we call the number $\frac{2N}{n}$ the 'degree' of 217 the graph (as a matter of fact $\frac{2N}{n}$ is the average degree of the vertices of G.) 218 If a graph G has the property that G has no subgraph having a larger degree 219 than G itself, we call G a balanced graph." 220

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Theorem 3.1 ([27]). Let $k \ge 2$ and $l \ (k-1 \le l \le {k \choose 2})$ be positive integers. 221 Let $\mathcal{B}_{k,l}$ denote an arbitrary not empty class of connected balanced graphs 222 consisting of k points and l edges. The threshold function for the property that 223 the random graph considered should contain at least one subgraph isomorphic 224 with some element of $\mathcal{B}_{k,l}$ is $n^{2-\frac{k}{l}}$.

Among special cases they mention trees, connected unicyclic graphs, 226 cycles, complete graphs and complete bipartite graphs all of which are 227 balanced. Over 20 years later, Bollobás [9] generalized this theorem to 228 arbitrary (not only balanced) graphs. He, however, used a rather complicated 229 method. In 1985, to a great surprise to all involved, Ruciński and Vince [73] 230 found out that the original proof of Erdős and Rényi which was based on 231 the second moment method can be easily adapted to cover all graphs as well. 232 We now state that result in the binomial model.

Theorem 3.2 ([9]). For an arbitrary graph G with at least one edge, 234

$$\lim_{n \to \infty} P(G \subset \Gamma_{n,p}) = \begin{cases} 0 & \text{if } p = o(n^{-1/m_G}) \\ 1 & \text{if } n^{-1/m_G} = o(p), \end{cases}$$
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where $m_G = \max_{H \subseteq G} d_H$ and $d_G = \frac{|E(G)|}{|V(G)|}$.

A crucial role in the Ruciński-Vince proof of Theorem 3.2 is played by 237 the quantity $\Phi_G = \min_{H \subseteq G} Exp(X_H)$. In fact, the inequalities 238

$$1 - \Phi_G \le P(G \not\subset \Gamma_{n,p}) \le c_1 / \Phi_G$$
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obtained in that proof have been strengthened to exponential bounds 240

$$e^{-c_2\Phi_G} \le P(G \not\subset \Gamma_{n,p}) \le e^{-c_3\Phi_G},$$
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where the L-H-R follows by the FKG inequality and the R-H-S is a special 242 case of a recent inequality from [42]. 243

As far as the asymptotic distributions of subgraph counts are concerned, 244 Erdős and Rényi treated in [ER60] only trees and cycles. For trees of order k 245 they established a limiting Poisson distribution on the threshold $N \sim cn^{\frac{k-2}{k-1}}$. 246 They observed that the same result holds for isolated trees, since in this 247 range almost surely all k-vertex trees are isolated (i.e., are components of the 248 random graph). They also found another Poisson threshold for isolated trees 249 at $N = \frac{1}{2k}n\log n + \frac{k-1}{2k}n\log\log n + cn + o(n)$, beyond which isolated trees die 250 out (swallowed by the giant component on its way to absorb all the vertices 251 of the random graph). They also established an asymptotic normality of the 252 number of isolated trees of order k (after suitable standardization) in the 253 whole range of N between the two thresholds. As observed by A. Barbour 254 in [5], the proof given by Erdős and Rényi was not correct and in the range 255 $N \sim cn, c \neq 1/2$, the standardization was not right. However, using another 256

method Barbour showed that indeed the asymptotic normality holds in the 257 entire range in question. For cycles and isolated cycles they established 258 a Poisson distribution (different in each case) at $N \sim cn$ and observed 259 that contrary to isolated trees, "...the probability that $\Gamma_{n,N}$ contains an 260 isolated cycle of order k never approaches 1." A similar result was proved for 261 connected unicyclic graphs. All these results were obtained by the method 262 of moments based on a fact from probability theory that for all distributions 263 which are uniquely determined by their moments (Poisson and normal are 264 such) the convergence of all moments of a sequence of random variables 265 to the moments of that distribution implies convergence in distribution 266 [8, Theorem 30.2]. Erdős and Rényi prove this fact as a lemma just for the 267 Poisson distribution, although they use it also for the normal distribution. 268 At the end of the paper, in a remark added in proof, they acknowledge that 269 N. V. Smirnov proved this lemma already in 1939. 270

They conclude their investigations of local properties of random graphs 271 with the comment: "Similar results can be proved for other types of subgraphs, 272 e.q., complete subgraphs of a given order. As however these results and their 273 proofs have the same pattern as those given above we do not dwell on the 274 subject any longer and pass to investigate global properties of the random 275 graph $\Gamma_{n,N}$." In 1979, K. Schürger, a former Ph.D. student of Erdős, proved 276 similar results for complete subgraphs [74] and a few years later Karoński [47] 277 extended them to so called k-trees, a common generalization of trees and 278 complete graphs. All these particular cases led to a general result for all 279 strictly balanced graphs. A graph is *strictly balanced* if every proper subgraph 280 has its degree strictly smaller than the graph itself. Let us denote $d_G = \frac{|E(\hat{G})|}{|V(G)|}$ 281 and recall that X_G is the number of copies of G in a random graph $\Gamma_{n,p}$. 282 The following result was proved independently in [9] and [48]. 283

Theorem 3.3 ([9, 48]). If G is a strictly balanced graph and $np^{d_G} \rightarrow c > 0$ 284 then X_G converges to the Poisson distribution with expectation $\frac{c^v}{aut(G)}$. 285

If a graph G is balanced but not strictly balanced then the limiting 286 distribution of X_G on the threshold, i.e. when $p = \Theta(n^{-1/d_G})$, becomes quite 287 involved. Although, in principle, as shown by Bollobás and Wierman [20], it 288 can be computed, there is no nice closed formula. For example, when G is a 289 disjoint union of 2 triangles then the limit distribution is that of the random 290 variable $\binom{Y}{2}$, where Y is Poisson. When G is the triangle with a pendant edge, 291 the limit is $\sum_{i=1}^{Z} Y_i$, where all random variables involved are independent 292 and Poisson. When G is the triangle with two pendant edges hanging at 293 the same vertex then X_G converges to the distribution of $\sum_{i=1}^{Z} \binom{Y_i}{2}$, where 294 again all random variables are independent Poisson. One more example: if 295 G is the triangle with a path of length 2 hanging at one of it vertices, then 296 the limit distribution is that of $\sum_{i=1}^{\Sigma_{j=1}^U W_j} Y_i$, where all random variables are 297 independent Poisson. We can only hope that so far the reader is convinced 298 that a pattern does indeed exist.

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If G is nonbalanced, then the expectation of X_G tends to infinity and one 300 has to normalize. It turns out that there is a nonrandom sequence $a_n(G) \to \infty$ 301 such that the asymptotic distribution of $\frac{X_G}{a_n(G)}$ coincides with that of X_H , 302 where H is the largest subgraph of G for which $d_H = m_G$. Clearly, H is 303 balanced and we are back to the balanced case. The sequence $a_n(G)$ is equal 304 to the expected number of extensions of a given copy of H to a copy of G in 305 the random graph $\Gamma_{n,p}$. For details see [71, page 292].

Beyond the threshold, i.e., when $np^{m_G} \to \infty$, X_G converges after 307 standardization to the standard normal distribution as long as $n^2(1-p) \to \infty$. 308 (For bigger $p X_G$ is either Poisson or degenerate, according to the formula 309 $X_G \sim {n \choose v} \frac{v!}{aut(G)} - c_n(G)Z$, where Z is the binomial random variable counting 310 edges in the complement of $\Gamma_{n,p}$ and $c_n(G)$ is the number of copies of G in K_n 311 containing a fixed edge. For details see [70].) This result was supplemented 312 by the rate of convergence in [7]. It was shown there that the total variation 313 distance between standardized X_G and the standard normal distribution can 314 be bounded by $O(\frac{1}{\sqrt{\Phi_G}})$ as long as $p \neq 1$ and by $O(\frac{1}{n\sqrt{1-p}})$ otherwise. Recall 315 that $\Phi_G \to \infty$ if and only if $np^{m_G} \to \infty$.

A variant of the small subgraph problem is one when we only count 317 induced subgraphs of $\Gamma_{n,p}$ which are isomorphic to G (induced copies). 318 Let Y_G count them. Then, denoting v = |V(G)| and l = |E(G)|, $Exp(Y_G) =$ 319 $Exp(X_G)(1-p)^{\binom{v}{2}-l}$, and as long as $p \to 0$ there is no substantial difference in 320 the limiting distribution of X_G and Y_G . For p constant, however, interesting 321 things may happen. First of all, in contrast to X_G , the variance of Y_G may 322 drop below the order of n^{2v-2} . It does so when $Exp(I|J_{12}) = Exp(I)$, i.e., 323 when $p = l/\binom{v}{2}$, where I is the indicator of the event that there is an induced 324 copy of G in $\Gamma_{n,p}$ on the vertex set $\{1,\ldots,v\}$ and J_{ij} is the indicator that 325 the edge ij is present in $\Gamma_{n,p}$. But if $Var(Y_G) = \Theta(n^{2\nu-3})$ then still Y_G is 326 asymptotically normal, and only when the variance drops further down to 327 the order of n^{2v-4} the distribution of standardized Y_G becomes nonnormal 328 (the convolution of normal and χ^2 distributions). It is a purely combinatorial 329 question when $Var(Y_G) = \Theta(n^{2\nu-4})$. For the higher terms to cancel out one 330 needs that $Exp(I|J_{12}, J_{13}, J_{23}) = Exp(I)$, or, equivalently, that in addition 331 to $p = l/\binom{v}{2}$, the proportion $t_3: t_2: t_1: t_0 = p^3: 3p^2q: 3pq^2: q^3$ is satisfied, 332 where t_i is the number of induced subgraphs of G isomorphic to the graph 333 with 3 vertices and i edges. For $p = \frac{1}{2}$, an example of a graph satisfying 334 these requirements is the wheel on 8 vertices, i.e. the graph obtained from 335 the 7-cycle by joining a new vertex to every vertex of the cycle. For some 336 time it was an open question if such abnormal cases take place for every 337 rational p. A positive answer to that puzzle is due to combined efforts of 338 Janson, Kratochvíl, Kärrman and Spencer [41, 45, 49]. 339

The random variables X_G and Y_G are examples of sums of random 340 variables with only few dependent summands. In particular, the summands 341 forming Y_G are dependent only if the sets corresponding to the indices 342 intersect (on at least 2 vertices, in fact). The reason is that the property of 343 380

the vertex set we are after depends only on the presence and absence of the 344 edges within the set. The situation changes when we move to the properties 345 depending also on the pairs with one endpoint in the set. Then all summands 346 are mutually dependent, but most just weakly. We have already encountered 347 such a case when studying the number of components of $\Gamma_{n,p}$ which are 348 isomorphic to a given graph G. Clearly this property requires that there is 349no edge with one endpoint in the set of vertices of a copy of G. Another 350 example of such "semi-induced" property is the notion of a maximal clique. 351 This is a complete subgraph not contained in any bigger complete subgraph 352 of a graph. For a vertex set to span a maximal clique one needs that no other 353 vertex is adjacent to all the vertices of the set. In [6] the limiting distribution 354 of the number of maximal k-cliques was investigated. It was proved that 355 for k > 2 there are two Poisson thresholds for the existence of maximal k- 356 cliques and the phase of asymptotic normality between them. Finally, there 357 are characteristics which lead to sums of random variables indexed by vertex 358 sets, which each depend on the presence or absence of all the edges in $\Gamma_{n,p}$. 359 An example of this is the number of copies of G disjoint from all other copies 360 of G in $\Gamma_{n,p}$. Here even the expectation is difficult to obtain, and the limiting 361normal distribution is still beyond ones reach. 362

4. Phase Transition

Sections 4–9 of [ER60] are devoted to global properties of random graphs. ³⁶⁴ The proofs follow the same pattern. First, the expectation of the quantity in ³⁶⁵ question is asymptotically evaluated. Then, using Markov's and Chebyshev's ³⁶⁶ inequality (the first and the second moment method, resp.) the asymptotics ³⁶⁷ of the quantities themselves are derived. As a summary of these results we ³⁶⁸ quote here how Erdős and Rényi characterize the process of the evolution ³⁶⁹ of a random graph in the paper presented to the International Statistical ³⁷⁰ Institute meeting in Tokyo in 1961 [ER61a]: ³⁷¹

"If n is fixed large positive integer and n is increasing from 1 to $\binom{n}{2}$, the 372 evolution of $\Gamma_{n,N}$ passes through five clearly distinguishable phases. These 373 phases correspond to ranges of growth of the number N of edges, these ranges 374 being defined in terms of the number n of vertices. 375

- Phase 1 corresponds to the range N(n) = o(n). For this phase it is 376 characteristic that $\Gamma_{n,N(n)}$ consists almost surely (i.e. with probability 377 tending to 1 as $n \to +\infty$) exclusively of components which are trees. (...) 378
- Phase 2 corresponds to the range $N(n) \sim cn$ with 0 < c < 1/2. (...) 379 In this range almost surely all components of $\Gamma_{n,N(n)}$ are either trees 380 or components consisting of an equal number of edges and vertices, i.e. 381 components containing exactly one cycle. (...) In this phase though not 382 all, but still almost all (i.e. n - o(n)) vertices belong to components which

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are trees. The mean number of components is n - N(n) + O(1), i.e. in 383 this range by adding a new edge the number of components decreases by 1, 384 except for the finite number of steps. 385

• Phase 3 corresponds to the range $N(n) \sim cn$ with $c \geq 1/2$. When N(n)386 passes the threshold n/2, the structure of $\Gamma_{n,N(n)}$ changes abruptly. As a 387 matter of fact this sudden change of the structure of $\Gamma_{n,N(n)}$ is the most 388 surprising fact discovered by the investigation of the evolution of random 389 graphs. While for $N(n) \sim cn$ with c < 1/2 the greatest component of 390 $\Gamma_{n,N(n)}$ is a tree and has (with probability tending to 1 as $n \to +\infty$) 391 approximately $\frac{1}{\alpha} \left(\log n - \frac{5}{2} \log \log n \right)$ vertices, where $\alpha = 2c - \log 2c$, for 392 $N(n) \sim n/2$ the greatest component has (with probability tending to 1 as 393 $n \rightarrow +\infty$) approximately $n^{2/3}$ vertices and has rather complex structure. 394 Moreover for $N(n) \sim cn$ with c > 1/2 the greatest component of $\Gamma_{n,N(n)}$ 395 has (with probability tending to 1 as $n \to +\infty$) approximately G(c)n 396 vertices, where 397

$$G(c) = 1 - \frac{1}{2c} \sum_{k=1}^{+\infty} \frac{k^{k-1}}{k!} \left(2ce^{-2c}\right)^k$$
398

(clearly G(1/2) = 0 and $\lim_{c \to +\infty} G(c) = 1$).

Except this "giant" component, the other components are all relatively 400 small, most of them being trees, the total number of vertices belonging to 401 components, which are trees being almost surely n(1 - G(c)) + o(n) for 402 $c \geq 1/2$. (...)

The evolution of $\Gamma_{n,N(n)}$ in Phase 3. may be characterized by that the 404 small components (most of which are trees) melt, each after another, into 405 the giant component, the smaller components having the larger chance of 406 "survival"; the survival time of a tree of order k which is present in $\Gamma_{n,N(n)}$ 407 with $N(n) \sim cn, c > 1/2$ is approximately exponentially distributed with 408 mean value n/2k.

- Phase 4 corresponds to the range $N(n) \sim cn \log n$ with $c \leq 1/2$. In this 410 phase the graph almost surely becomes connected. (...) 411
- Phase 5 consists of range $N(n) \sim (n \log n)\omega(n)$ where $\omega(n) \to +\infty$. 412 In this range the whole graph is not only almost surely connected, but the 413 orders of all points are almost surely asymptotically equal. Thus the graph 414 becomes in this phase "asymptotically regular". "415

Erdős and Rényi in their fundamental paper [ER60] gave a fairly complete 416 "big picture" of the evolution of a random graphs. However many fascinating 417 questions were left unanswered. For example, how did the giant component 418 grow so rapidly, what is the nature of the "double jump" of its size: from 419 $O(\log n)$ when c < 1/2 to $\Theta(n^{2/3})$ when c = 1/2 and finally being of the 420 order of n when c > 1/2? 421

381

382Often we say that a random graph goes through the phase transition at 422

Author's Proof

c = 1/2 due to an obvious resemblance of this period of its evolution to the 423 physical phenomena of changing the state, for example, from liquid to solid. 424 Here a random graph changes abruptly its state from a loose collection of 425 small components being trees and unicyclic to solid single giant component 426 dominating its structure. 427

The critical moment of the phase transition was unresolved until the 428 milestone paper of Béla Bollobás [13] who revealed the mechanism of the 429 formation of the giant component. He also focused the attention, for the first 430 time, on the nature of the phase transition phenomena, investigating this 431 critical moment of the evolution and looking at the beginning of so called 432 supercritical phase. He asked what is the typical structure of a random graph 433 $\Gamma_{n,N}$ when $N(n) = \frac{1}{2}n + s$, where s = o(n). In particular he proved that 434 the largest component is almost surely unique once $s \ge 2(\log n)^{1/2}n^{2/3}$ and 435 its size $L_1(\Gamma_{n,N})$ is approximately 4s while the size of the second largest 436 component $L_2(\Gamma_{n,N})$ is much smaller. 437

Bollobás gave a good lead to what we might consider as the proper 438 magnification if we want to get undistorted picture of the phase transition 439 while looking at the neighborhood of the "critical point" n/2. Due to later 440 results of Luczak [58], combined with those of Kolchin [51], we know that the 441 correct parametrization is 442

$$N(n) = \frac{1}{2}n + \lambda n^{2/3}.$$
 443

When $\lambda \to -\infty$ then $\Gamma_{n,N}$ consists of many components of the same 444 size as the largest one, which is still very small and consists roughly of 445 $\frac{n^2}{2c^2}\log(s^3/n^2)$ vertices, and the large components are unable to "swallow" 446 each other and therefore are forced to hunt for smaller query. Hence large 447 components grow absorbing only small ones and no clear favorite to win the 448 race for the giant emerges. As the number of edges N(n) increases, the number 449 of contestants decreases. When $\lambda = constant < 0$ the probability that two 450 specified large components will form a new component is bounded away from 451 zero, but still too small to ensure the creation of unique giant component. At 452 the same time, a big gap between the orders of large and small components 453 arises which prevents the creation of new large components from the small 454 ones. Next, as soon as $\lambda \to \infty$, all large components almost "instantly" merge 455 together and a unique large component emerges. This component is still not 456 giant, it has barely over $n^{2/3}$ vertices, but it will continue to absorb other 457 components, first the largest ones, rapidly becoming giant. 458

The next result of Luczak [58] gives a clear picture of the sizes $L_i(\Gamma_{n,N})$ 459 of the ith largest components during the phase transition of $\Gamma_{n,N}$. Here and 460 throughout the paper the abbreviation a.s. stands for 'almost surely', a phrase 461 whose precise meaning was explained in the description of Phase 1. above. 462

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Theorem 4.1 ([58]). Let k be natural number and $sn^{-2/3} \rightarrow \infty$ but s = 463 o(n).

(i) If
$$N = n/2 - s$$
 then for every $i = 1, 2, ..., k$ and every real r

$$\lim_{n \to \infty} P\left(L_i(\Gamma_{n,N}) < \frac{n^2}{2s^2} \left(\log \frac{s^3}{n^2} - \frac{5}{2}\log \log \frac{s^3}{n^2} + r\right)\right) = \sum_{j=0}^{i-1} \frac{\lambda^j}{j!} e^{-\lambda}, \quad 466$$

where $\lambda = \lambda(r) = 2/\sqrt{\pi}e^{-r}$.

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Moreover, a.s. the *i*th largest component of $\Gamma_{n,N}$ is a tree for i = 4681,2,..., k and $\Gamma_{n,N}$ contains no component with more edges than vertices. 469 (ii) Let N = n/2 + s and let s' be the unique positive solution of the equation 470

$$\left(1 - \frac{2s'}{n}\right)e^{\frac{2s'}{n}} = \left(1 + \frac{2s'}{n}\right)e^{-\frac{2s'}{n}}.$$
471
472

Then a.s.

$$\left|L_1(\Gamma_{n,N}) - \frac{2(s+s')n}{n+2s}\right| < \omega(n)\frac{n}{\sqrt{s}}$$

$$473$$

and so

$$|L_1(\Gamma_{n,N}) - 4s| < \omega(n)\frac{n}{\sqrt{s}} + O\left(\frac{s^2}{n}\right).$$
475

Moreover, for every i = 2, ..., k and every real r

$$\lim_{n \to \infty} P\left(L_i(\Gamma_{n,N}) < \frac{n^2}{2s^2} \left(\log \frac{s^3}{n^2} - \frac{5}{2}\log\log \frac{s^3}{n^2} + r\right)\right) = \sum_{j=0}^{i-1} \frac{\lambda^j}{j!} e^{-\lambda}, \quad 477$$

where $\lambda = \lambda(r) = 2/\sqrt{\pi}e^{-r}$.

Furthermore a.s. the ith largest component of $\Gamma_{n,N}$, i = 2, 3, ..., k, 479 is a tree and no component of $\Gamma_{n,N}$, except for the largest one, contains 480 more edges than vertices. 481

To study the critical "interval" when the phase transition takes place, 482 i.e., when $N(n) = \frac{1}{2}n + \lambda n^{2/3}$ and $\lambda \to \mp \infty$, requires very sophisticated and 483 delicate tools. Janson, Knuth, Luczak and Pittel in their extensive, almost 484 140 pages long, study [40] applied machinery of generating functions with 485 great success. They were able to analyze the structure of evolving graphs 486 (and multigraphs) when edges are added one at a time and at random, with 487 great precision, mainly looking and so called excess and deficiency of a graph. 488 To give the reader a taste of their results let us quote the following theorem. 489 **Theorem 4.2** ([40]). The probability that a random graph or multigraph 490 with n vertices and $\frac{1}{2}n + O(n^{1/3})$ edges has exactly r bicyclic components 491 (i.e., components with exactly two cycles), and no components of higher cyclic 492 order, is 493

$$\left(\frac{5}{18}\right)^r \sqrt{\frac{2}{3}} \frac{1}{(2r)!} + O(n^{-1/3}).$$
494

They also study the following fascinating problem: What is the 495 probability that the component which during the evolution becomes the first 496 "complex" component (i.e., the first component with more than one cycle) 497 is the only complex component which emerges during the whole process? 498 So they ask what is the probability that the first bicyclic component is the 499 "seed" for the giant one. They prove that it happens quite often indeed. 500

Theorem 4.3 ([40]). The probability that an evolving graph or multigraph 501 on n vertices never has more than one complex component throughout its 502 evolution approaches $\frac{5\pi}{18} \approx 0.8727$ as $n \to \infty$. 503

5. Planarity and Chromatic Number

In a paper of such an enormous length one can likely find less rigorous claims. 505 One of such things happened in the paper [ER60] in relation to the question 506 when a random graph $\Gamma_{n,N}$ is planar. 507

Since trees and components with exactly one cycle are planar, Erdős and 508 Rényi easily deduced from their findings about early stages of the evolution 509 of a random graph, that when c < 1/2 then the probability that $\Gamma_{n,N}$ is 510 planar tends to 1. Now, to support the claim that when c passes 1/2 the 511 graph becomes non-planar they used the argument that $\Gamma_{n,N}$ contains an 512 induced cycle with d diagonals. Although their claim (Theorem 8a on page 513 51) regarding the distribution of the number of such cycles is incorrect, as 514 it was pointed out later by Luczak and Wierman [63], their intuition was 515 perfect and the following result is indeed true.

Theorem 5.1 ([63]). Let us suppose that $N \sim cn$. If c < 1/2 the probability 517 that the graph $\Gamma_{n,N}$ is planar is tending to 1 while for c > 1/2 this probability 518 tends to 0. 519

Such a behavior of a random graph shows the fundamental difference 520 in its typical structure before and after the phase transition. Now, thanks 521 to the contribution of Luczak, Pittel and Wierman [64], we have more 522 detailed knowledge about planarity of a random graph, also during the phase 523 transition. 524

Theorem 5.2 ([64]). Let $\epsilon = \epsilon(n) \to 0$ as $n \to \infty$. Then $\Gamma_{n,p}$ is:

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Author's Proof

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- (i) a.s. planar, when $p = (1 \epsilon)/n, \epsilon^3 n \to \infty$; 526
- (ii) Planar with probability tending to $a(\lambda), 0 < a(\lambda) < 1$, as $n \to \infty$, when 527 $p = (1 + \epsilon)/n$, where $\epsilon^3 n \to \lambda$ and $-\infty < \lambda < \infty$ is a constant: 528 529
- (iii) a.s. non-planar, when $p = (1 + \epsilon)/n$, $\epsilon^3 n \to \infty$.

In the final section of the paper [ER60] Erdős and Rényi collected 530 unsolved problems. One of them is closely related to planarity: Another in- 531 teresting question is: what is the threshold for the appearance of a "topological 532 complete graph of order k". i.e., of k points such that any two of them can 533 be connected by a path and these paths do not intersect. For k > 4 we do 534 not know the solution. The solution was found many years later by Ajtai, 535 Kómlos, and Szemerédi [2]. 536

Another problem mentioned there turned out to be one of the central 537 and most challenging questions of the theory. Erdős and Rényi asked "what 538 will be the chromatic number of $\Gamma_{n,N}$?" What they knew then about this 539 important graph invariant was limited to facts which can be deduced from 540 general results regarding the evolutionary process. Here is what they were 541 able to conclude : "Clearly every tree can be colored by 2 colors, and thus 542 by Theorem 4a almost surely $Ch(\Gamma_{n,N}) = 2$ if N(n) = o(n). As however 543 the chromatic number of a graph having an equal number of vertices and 544 edges is equal to 2 or 3 according whether the only cycle contained in such 545 graph is of even or odd order, it follows from Theorem 5e that almost surely 546 $Ch(\Gamma_{n,N}) \leq 3$ for $N(n) \sim nc$ with c < 1/2. For $N(n) \sim n/2$ we have 547 almost surely $Ch(\Gamma_{n,N}) \geq 3$. As a matter of fact, in the same way, as we 548 proved Theorem 5b, one can prove that $\Gamma_{n,N}$ contains for $N(n) \sim n/2$ almost 549 surely a cycle of odd order. It is an open problem how large $Ch(\Gamma_{n,N})$ is for 550 $N(n) \sim n/2$ with c > 1/2." 551

This question remained open for next 30 years, and was answered, for 552 large c, by Luczak in [57]. He proved that the chromatic number $\chi(\Gamma_{n,n})$ 553 behaves as follows. 554

Theorem 5.3 ([57]). Let np = c and $\epsilon > 0$ be fixed. Suppose $c_{\epsilon} \leq c + o(n)$ 555 for sufficiently large constant c_{ϵ} . Then 556

$$P\left(\frac{c}{2\log c} < \chi(\Gamma_{n,p}) < (1+\epsilon)\frac{c}{2\log c}\right) \to 1 \quad as \quad n \to \infty.$$
 557

Although the original question was posed for sparse random graphs 558 the ideas leading to the proof came from investigations of the chromatic 559 number of dense random graphs. The first step toward the solution was 560 made by Matula [66, 67] and Bollobás and Erdős [16] who discovered high 561 concentration of the size of the largest independent set in $\Gamma_{n,p}$ around $2\log_b n$, 562 where b = 1/(1-p) and edge probability p is a constant. It suggested that the 563 respective lower bound for $\chi(\Gamma_{n,p})$ should be $n/(2\log_b n)$. Only a few years 564 later, Grimmett and McDiarmid published a paper [38] in which they showed 565 that a greedy algorithm, which assigns colors to vertices of a random graph 566 sequentially, in such a way that a vertex gets the first available color, needs, 567 with high probability, approximately $n/\log_b n$ colors to produce a proper 568 coloring of $\Gamma_{n,p}$. It established an upper bound for the chromatic number 569 of dense random graph, twice as large as the lower bound. Grimmett and 570 McDiarmid conjectured that the lower bound sets, in fact, the correct order of 571 magnitude for $\chi(\Gamma_{n,p})$. The right tool to settle this conjecture was delivered 572 by Shamir and Spencer [76]. They proved that the chromatic number of 573 $\Gamma_{n,p}$ is sharply concentrated in an interval of length of order $n^{1/2}$ but, what 574 perhaps was more important then their result itself, they introduced to the 575 theory of random graphs a new powerful technique based on concentration 576 measure of martingales, known in the probabilistic literature as Hoeffding-577 Azuma inequality. But it was Béla Bollobás who showed how the potential of 578 martingale approach can be utilized to solve long standing conjecture. In his 579 paper [15] he proved the following theorem. 580

Theorem 5.4 ([15]). Let 0 be fixed and <math>b = 1/(1-p). Then for 581 every $\epsilon > 0$ 582

$$P(\frac{n}{2\log_b n} < \chi(\Gamma_{n,p}) < (1+\epsilon)\frac{n}{2\log_b n}) \to 1 \quad as \quad n \to \infty.$$
583

Later on Matula and Kucera [68] gave an alternative proof of the above 584 theorem, using the second moment and "expose and merge" algorithmic 585 approach. Luczak's proof of Theorem 5.3 is in fact an ingenious blend of 586 the martingale and "expose and merge" techniques. 587

The chromatic number of a random graph is a random variable, the 588 distribution of which should be highly concentrated. It is easy to notice 589 (see above) that if $p = o(n^{-1})$ then $\chi(\Gamma_{n,p})$ is 2 (not counting the case when 590 the edge probability is of the order smaller then n^{-2} and therefore, with high 591 probability the graph is empty). One can also show that when $p \sim cn^{-1}$, 592 O < c < 1 then $P(\chi(\Gamma_{n,p}) = 2) \rightarrow a$ and $P(\chi(\Gamma_{n,p}) = 3) \rightarrow 1 - a$, where 593 $a = e^{c/2}((1-c)/(1+c))^{1/4}$. The last probabilities are simply the same as the 594 probabilities that $\Gamma_{n,p}$ has or does not have an odd cycle. Such a behavior of 595 a random variable χ has been confirmed, for small edge probabilities only, by 596 Luczak. He proved in [61] that if $p < n^{-5/6-\epsilon}$ then the chromatic number, as 597 expected, takes on at most two values.

6. Asymmetric Graphs

Another interesting topic originated from a joint paper by Erdős and Rényi $_{600}$ in the peak of their cooperation in early 1960s [ER63]. Here is how they $_{601}$ describe their goals: "We shall call (...) a graph symmetric, if there exists $_{602}$ a non-identical permutation of its vertices, which leaves the graph invariant. $_{603}$ By other words, a graph is called symmetric if the group of its automorphisms $_{604}$ has degree greater than 1. A graph which is not symmetric will be called $_{605}$

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asymmetric. The degree of symmetry of a symmetric graph is evidently 606 measured by the degree of its group of automorphisms. The question which 607 led us to the results contained in the present paper is the following: how can 608 we measure the degree of asymmetry of an asymmetric graph?" 609

They answer the last question in what follows: "Evidently any asymmetric 610 graph can be made symmetric by deleting certain of its edges and by adding 611 certain new edges connecting its vertices. We shall call such a transformation 612 of the graph its symmetrization. For each symmetrization of the graph let us 613 take the sum of the number of deleted edges – say r – and the number of new 614 edges – say s –; it is reasonable to define the degree of asymmetry A[G] of 615 a graph G, as the minimum of r + s where the minimum is taken over all 616 possible symmetrizations of the graph G. (...) The question arises: how large 617 can be the degree of asymmetry of a graph of order n (i.e., a graph which has 618 n vertices)? We shall denote by A(n) the maximum of A[G] for all graphs G 619 of order n(n = 2, 3, ...)."

They first notice that A(2) = A(3) = A(4) = A(5) = 0 while A(6) = 6211. In general, a rather straightforward deterministic argument leads to the 622 following result. 623

Theorem 6.1 ([30]).

$$A(n) \le \left\lfloor \frac{n-1}{2} \right\rfloor.$$
625

To find the lower bound for A(n) Erdős and Rényi use a non-constructive 626 argument, i.e., they show via the probabilistic method that there exists a 627 certain graph on n vertices with the degree of asymmetry at least $n(1-\epsilon)/2$, 628 $0 < \epsilon < 1$.

Theorem 6.2 ([30]). Let us choose at random a graph Γ having n given 630 vertices so that all possible $2^{\binom{n}{2}}$ graphs should have the same probability 631 to be chosen. Let $\epsilon > 0$ be arbitrary. Let $P_n(\epsilon)$ denote the probability that 632 by changing not more than $\frac{n(1-\epsilon)}{2}$ edges of Γ it can be transformed into a 633 symmetric graph. Then we have 634

$$\lim_{n \to \infty} P_n(\epsilon) = 0.$$
 635

Corollary 6.1. For any ϵ with $0 < \epsilon < 1$ there exists an integer $n_0(\epsilon)$ 636 depending only on ϵ , such that for every $n > n_0(\epsilon)$ there exists a graph G of 637 order n with $A[G] > n(1-\epsilon)/2$.

Indeed, for large n, Theorem 6.2 shows that almost every graph is a 639 counterexample to the hypothesis that its symmetrization is possible with 640 less than $\frac{n}{2}(1-o(1))$ edges. 641

387



Hence, if we combine Theorem 6.1 and Corollary 6.1 we see that

$$\lim_{n \to \infty} \frac{A(n)}{n} = \frac{1}{2}.$$
643

After showing that almost all labeled simple graphs are asymmetric, Erdős 644 and Rényi turned their attention to graphs with a prescribed number of 645 edges. First they noticed that since almost every tree has a cherry, i.e., a pair 646 of pendant vertices adjacent to a common neighbor, therefore almost every 647 tree on n vertices is symmetric. Furthermore they proved that any connected 648 graph of order n having n edges is either symmetric or its asymmetry is one 649 and gave the following bound. 650

Theorem 6.3 ([30]). If a graph G of order n has $N = \lambda n$ edges $(0 < \lambda < 651 (n-1)/2)$ then 652

$$A[G] \le 4\lambda \left(1 - \frac{2\lambda}{n-1}\right).$$
⁶⁵³

Erdős and Rényi went further in their investigations. Let us quote a 654 few more lines from their paper [ER63]. "Another interesting question is 655 to investigate the asymmetry or symmetry of a graph for which not only the 656 number of vertices but also the number of edges N is fixed, and to ask that 657 if we choose one of these graphs at random, what is the probability of its 658 being asymmetric. We have solved this question too, and have shown that if 659 $N = \frac{n}{2}(\log n + \omega(n))$, where $\omega(n)$ tends arbitrarily slowly to $+\infty$ for $n \to +\infty$, 660 then the probability that a graph with n vertices and N edges chosen at random 661 (so that any such graph has the same probability $\binom{\binom{n}{2}}{2}^{-1}$ to be chosen) should 662 be asymmetric, tends to 1 for $n \to +\infty$. This and some further results will 663 be published in another forthcoming paper."

Unfortunately the announced paper has never been published! Several 665 years later this problem and the analogous one for unlabeled graphs was 666 attacked again by Wright [79]. 667

Consider graphs $\Gamma_{n,N}$ and $U_{n,N}$ picked at random from the families of 668 all labeled and unlabeled graphs on n vertices and with N = N(n) edges, 669 respectively. Here is the result of Wright. 670

Theorem 6.4 ([79]). If $\omega(n) = (2N(n)/n) - \log n \to \infty$ then $\Gamma_{n,N}$ and $U_{n,N}$ 671 are almost surely asymmetric while when $\omega(n) \le 0$ then they are almost surely 672 symmetric. 673

Later Luczak [56] gave precise results about the structure of the 674 automorphism group $Aut(\Gamma_{n,N})$ of a random graph $\Gamma_{n,N}$. He studied the 675 symmetry of the largest component $L_1(n,N)$ of this random graph. What 676 he found was that when $N(n) = \frac{1}{2}n\alpha(n)$ then there exists a constant d such

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that for $\alpha(n) \ge d$ almost surely $Aut(L_1(n, N))$ is isomorphic to some product 677 of symmetric groups. From this result he was able to deduce the following 678 strengthening of the "labelled" part of Theorem 6.4. 679

Theorem 6.5 ([56]). Let $N = \frac{n}{2}(\log n + \omega(n)).$ 680

(i) If $\omega(n) \to -\infty$ then $|Aut(\Gamma_{n,N})| \to \infty$ a.s. (ii) If $\omega(n) \to c$ then

$$\lim_{n \to \infty} P(|Aut(\Gamma_{n,N})| = 1) = e^{\lambda} (1+\lambda)$$
$$\lim_{n \to \infty} P(|Aut(\Gamma_{n,N})| = k!) = \frac{\lambda^k}{k!} e^{-\lambda}$$

for $k = 2, 3, ..., where \lambda = e^{-c}$ and c is a constant. (iii) If $\omega(n) \to \infty$ then $|Aut(\Gamma_{n,N})| = 1$ a.s.

7. Perfect Matchings

The last three papers Erdős and Rényi wrote on the subject of random graphs 686 were devoted to the existence of 1-factors. In [ER64] and [ER68] they coped 687 with the relatively easier case of random bipartite graphs. In both papers 688 they consequently emphasized the matrix terminology. "In the present paper 689 we deal with certain random 0-1 matrices. Let $\mathcal{M}(n,N)$ denote the set of 690 all n by n square matrices among the elements of which there are exactly N $_{691}$ elements $(n \leq N \leq n^2)$ equal to 1, all the other elements are equal to 0. The 692 set $\mathcal{M}(n, N)$ contains clearly $\binom{n^2}{N}$ such matrices; we consider a matrix M 693 chosen at random from the set $\mathcal{M}(n, N)$, so that each element of $\mathcal{M}(n, N)$ 694 has the same probability $\binom{n^2}{N}^{-1}$ to be chosen. We ask how large N has to 695 be, for a given large value of n, in order that the permanent of the random 696 matrix M should be different from zero with probability $\geq \alpha$, where $0 < \alpha < 1$. 697 (\dots) A second way to formulate the problem is as follows: we shall say that 698 two elements of a matrix are in independent position if they are not in the 699 same row and not in the same column. Now our question is to determine 700 the probability that the random matrix M should contain n elements which 701 are all equal to 1 and pairwise in independent position." 702

The result they prove resembles that for the connectedness (compare 703 Theorem 2.1). 704

Theorem 7.1 ([31]). Let P(n, N) denote the probability of the event that 705 the permanent of the random matrix M is positive. Then if 706

$$N(n) = n\log n + cn + o(n)$$
⁷⁰⁷

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where c is any real constant, we have

$$\lim_{n \to \infty} P(n, N(n)) = e^{-2e^{-c}}.$$
710

Finally, they also mention graphs: "This result can be interpreted also in 711 the following way, in terms of graph theory. Let $\Gamma_{n,N}$ be a bichromatic random 712 graph containing n red and n blue vertices, and N edges which are chosen at 713 random among the n² possible edges connecting two vertices having different 714 color (so that each of the $\binom{n^2}{N}$ possible choices has the same probability). 715 Then P(n,N) is equal to the probability that the random graph $\Gamma_{n,N}$ should 716 contain a factor of degree 1, i.e., $\Gamma_{n,N}$ should have a subgraph which contains 717 all vertices of $\Gamma_{n,N}$ and n disjoint edges, i.e., n edges which have no common 718 endpoint." (They seem not to use the name 'perfect matching' at all.) 719

As far as the proof is concerned, "Besides elementary combinatorial and 720 probabilistic arguments similar to that used by us in our previous work on 721 random graphs (...) our main tool in proving our results is the well-known 722 theorem of D. König, which is nowadays well known in the theory of linear 723 programming, according to which if M is an n by n matrix, every element 724 of which is either 0 or 1, then the minimal number of lines (i.e., rows or 725 columns) which contain all the 1-s, is equal to the maximal number of 1-s in 726 independent position. As a matter of fact, for our purposes we need only the 727 special case of this theorem, proved already by Frobenius (1917), concerning 728 the case when the maximal number of ones in independent positions is equal 729 to n (...). According to the theorem of Frobenius-König 1 - P(n, N) is equal 730 to the probability that there exists a number k such that there can be found k 731 rows and n-k-1 columns of M which contain all the ones $(0 \le k \le n-1)$." 732 The rest of the proof is devoted to showing that this is very unlikely for 733 N(n) given. It is interesting to notice that Erdős and Rényi never mention 734 Hall's theorem, which is equivalent to Frobenius but far more popular in 735 combinatorics nowadays. 736

The 1968 paper is a straightforward extension of the 1964 result, where 737 it is shown that setting 738

$$N(n) = n\log n + (r-1)n\log\log n + n\omega(n)$$
739

where $\omega(n)$ tends arbitrarily slowly to infinity then almost surely the 740 bichromatic random graph contains r disjoint 1-factors. The only new element 741 of the proof is the observation that if there are no r disjoint 1-factors then 742 there is a way to delete some edges so that no vertex looses more than r-1 743 from its degree and the resulting subgraph contains no 1-factor at all. Then 744 again the theorem of Frobenius is used. 745

The most involved of the three papers about 1-factors is that from 746 1966, where an ordinary (not bichromatic) random graph $\Gamma_{n,N}$ is considered. 747 The reason is that the theorem of Tutte describing the structure of graphs 748 which admit 1-factors is more complex than its counterpart in the bipartite 749

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case. "It should be added that the problem investigated in the present paper 750 is much more difficult than the corresponding problem for even graphs solved 751 in [5]. Thus for instance in [5] we made use of the well known theorem of 752 D. König; the corresponding tool in the present paper is the much deeper 753 theorem of Tutte mentioned above." ([5] = [ER 64]) 754

The result of that paper says that the threshold for containing a 755 1-factor coincides with that for disappearance of isolated vertices, and thus 756 also with that for connectivity (see Theorem 2.1). The proof is long and 757 tedious and involves a weaker version of Tutte's theorem ignoring the parity 758 of components. 759

Erdős and Rényi make also the following claim: "If $N = \frac{1}{2}n \log n + O(n)$, 760 as mentioned above, with probability near to 1, $\Gamma_{n,N}$ consists of a connected 761 component and a certain number of isolated points. With the same method 762 (...) one can prove that if the connected component of $\Gamma_{n,N}$ consists of an 763 even number of points, it has with probability near 1 a factor of degree one. 764 As the proof of this result is almost the same (...) we do not go into the 765 details."

The above mentioned result was proved (in a strengthened form) by 767 Bollobás and Thomason [18]. In order to quote that result let us extend 768 the notion of a perfect matching by saying that a graph satisfies property 769 \mathcal{PM} if there is a matching covering all but at most one of the nonisolated 770 vertices. It is known that, switching to the binomial model, as soon as 771 $2np - \log n - \log \log n \to \infty$, there are only isolated vertices outside the 772 giant component. However, the main obstacle for the property \mathcal{PM} is the 773 presence of a pair (at least two such pairs when the number of nonisolates is 774 odd) of vertices of degree 1 adjacent to the same vertex (called, as we already 775 mentioned, 'a cherry'). The expected number of cherries is 776

$$3\binom{n}{3}p^2(1-p)^{2(n-3)} < n^3p^2e^{-2np+6p} = o(1)$$
⁷⁷⁷

if $2np - \log n - 2\log \log n \to \infty$. Again, a trivial necessary condition becomes 778 almost surely sufficient. 779

Theorem 7.2 ([18]). Let
$$y_n = 2np - \log n - 2\log \log n \to \infty$$
. Then 780

$$P(\Gamma_{n,p} \in \mathcal{PM}) \to \begin{cases} 0 & \text{if } y_n \to -\infty \\ e^{-\frac{1}{8}e^{-c}} & \text{if } y_n \to c \\ 1 & \text{if } y_n \to \infty. \end{cases}$$
781

The proof, again, was based on Tutte's theorem. Years later Luczak and 782 Ruciński proposed an alternative approach, via Hall's Theorem, invented 783 in [65] to attack a more general question. For a given graph G, a perfect 784 G-matching of a graph is a spanning subgraph which is a disjoint union of 785 copies of G. For $G = K_2$ this is the ordinary notion of a 1-factor. 786 In [65] it was shown that for every nontrivial tree T, the threshold is the 787 same as that for disappearance of isolated vertices. 788

Theorem 7.3 ([65]). For every tree T on t vertices and with at least one 789 edge, assuming n is divisible by t, 790

$$P(\Gamma_{n,p} \text{ has a perfect } T \text{-matching }) \to \begin{cases} 0 & \text{if } np - \log n \to -\infty \\ e^{-e^{-c}} & \text{if } np - \log n \to c \\ 1 & \text{if } np - \log n \to \infty. \end{cases}$$

The threshold for arbitrary G is not known in general. Some partial results 792 are contained in [4] and [72]. 793

Coming back to the original papers of Erdős and Rényi, the last of them 794 is concluded by the following problem: "does a random graph $\Gamma_{n,N}$ where n 795 is even and 796

$$N = \frac{1}{2}n\log n + \frac{r-1}{2}n\log\log n + \omega(n)n$$
797

where $\omega(n) \to \infty$, contain at least r disjoint factors of degree one with 798 probability tending to 1 for $n \to \infty$?" 799

Shamir and Upfal [77] answered this question in the positive. Given a 800 map f of V(G) into the set of non-negative integers, define an f-factor of G 801 as a spanning subgraph of G in which the degree of vertex x is f(x). 802

Theorem 7.4 ([77]). *If*

$$\rho = \frac{1}{n} (\log n + (r-1)\log\log n + \omega(n)), \qquad 804$$

 $r \geq 1$, $\lim_{n \to \infty} \omega(n) = \infty$ and $1 \leq f(x_i) \leq r$, $\sum_{i=1}^n f(x_i)$ even, then $\Gamma_{n,p}$ has 805 an f-factor, almost surely.

Although *f*-factors are characterized by Tutte's theorem, Shamir and 807 Upfal chose an alternative approach using an algorithmic technique (in-808 troduced to random graphs by Pósa) of augmentation of sub-factors by 809 alternating paths. In fact, the answer to the last question of Erdős and 810 Rényi does not follow directly from the above result (not every *r*-factor 811 has a 1-factorization) but from the proof. In 1985 Bollobás and Frieze [17] 812 strengthened this answer by proving that almost surely in the random graph 813 process of adding edges one by one, as soon as the minimum degree becomes 814 r, there are $\lfloor r/2 \rfloor$ disjoint hamiltonian cycles plus a disjoint perfect matching 815 if r is odd.

The next problem we would like to mention cannot be directly attributed 817 to Erdős and Rényi. Here is how Erdős describes their omission [3, Ap- 818 pendix B]. "When Rényi and I developed our theory of random graphs, we 819 thought of extending our study for hypergraphs. We mistakenly thought that 820

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all (or most) of the extensions would be routine and we completely overlooked s21 the following beautiful question of Shamir. (...) Shamir asked how many s22 triples must one choose on 3n elements so that with probability bounded away s23 from zero one should get n vertex disjoint triples. Shamir proved that $n^{3/2}$ s24 triples suffice, but the truth may very well be $n^{1+\epsilon}$ or even $cn \log n$. The reason s25 for the difficulty is that Tutte's theorem seem to have no analogy for triple s26 systems or more generally for hypergraphs." The result mentioned by Erdős s27 belongs, in fact, to J. Schmidt-Pruzan and E. Shamir [75]. In 1995, Frieze s28 and Janson in [35] pushed the bound down to $n^{4/3}$.

Fortunately, Erdős and Rényi did not overlook some other important s30 problems which stimulated the research in the theory of random graphs 831 over the years. One such problem was the threshold for existence of a 632 Hamiltonian cycle in a random graph. They, in fact, asked only: for what order 633 of magnitude of N(n) has $\Gamma_{n,N(n)}$ with probability tending to 1 a Hamiltons34 line (i.e., a path which passes through all vertices). This problem was first 635 tried by Pósa [69] and Korshunov [55] and finally solved by Kómlos and 636 Szemerédi [54] and, in a stronger form, by Bollobás [12]. They proved that 637 the threshold for Hamiltonian cycle coincides with that of disappearance of 638 all vertices of degree 0 and 1.

8. Update for the Second Edition

We wrote this paper back in 1995. In this second edition of the volume we 841 decided to leave the original text intact except for a few obvious corrections 842 and the proofs of Theorems 3.2 and 3.3 which have been deleted entirely. 843 However, several new developments have occurred afterward. Here we would 844 like to mention some of them along with a couple of earlier results omitted 845 in the first edition. Needles to say, our choice is quite subjective. For more 846 thorough treatment of random graphs we refer the reader to the monograph 847 [43] published in 2000. 848

In relation to connectivity, one should note that an old result of 849 Luczak [60] states that the k-core of a random graph, for p large enough, 850 is a.s. empty or k-connected. It implies that $\Gamma_{n,p}$ is a.s. $c(\Gamma_{n,p})$ -connected for 851 the ranges of N larger than those in Theorem 2.2. 852

In the domain of small subgraphs of random graphs there has been an e_{53} intense study of the so called *upper tail* of the random variable X_G counting e_{54} copies of a given graph G in $\Gamma_{n,p}$. As far as the *lower tail* is concerned, whose e_{55} special case is the probability P(X = 0) discussed briefly after Theorem 3.2, e_{56} the asymptotic order of magnitude of the logarithm of $P(X \leq (1 - \epsilon)EX)$ est has been determined in [39] to be $-\Phi_G$. The exponent in the upper tail, e_{55} be determined. In [44] general lower and upper bounds were obtained which

differ only by a logarithmic factor. Very recently DeMarco and Kahn [21] have 860 found the right threshold for cliques and formulated the "right" conjecture 861 for the general case. 862

In Sect. 5, the threshold for topological cliques found in [2] has been 863 sharpened (see [62], a remark after Corollary 18). A significant result about 864 the chromatic number of a random graph appeared in [1]. Achlioptas and 865 Naor found therein an explicit two-point limiting distribution of $\chi(\Gamma_{n,p})$, 866 where p = d/n, for every d > 0, strengthening a theorem from [61] mentioned 867 at the end of Sect. 5. 868

The most acclaimed result in random graph theory which appeared after 869 1995 is, without doubt, a solution to the celebrated Shamir problem posed 870 in Sect. 7. After some initial attempts (Krivelevich [52, 53] and Kim [50]), in 871 2008 Johansson, Kahn, and Vu [46] published a complete solution to both, the 872 hypergraph Shamir problem and to its random graph counterpart (triangle- 873 factors), receiving for their achievement the prestigious Fulkerson Prize. Quite 874 recently in a series of papers, Dudek, Frieze, Loh, and Speiss [22-24, 34] 875 obtained thresholds for the hamiltonicity of random uniform hypergraphs. In 876 the hardest case of so called loose Hamilton cycles they incorporated in their 877 proofs the result on perfect matchings from [46]. 878

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